WADS-DSN 2004 Invited Talk

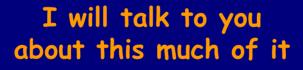
From Dependable Architectures To Dependable Systems

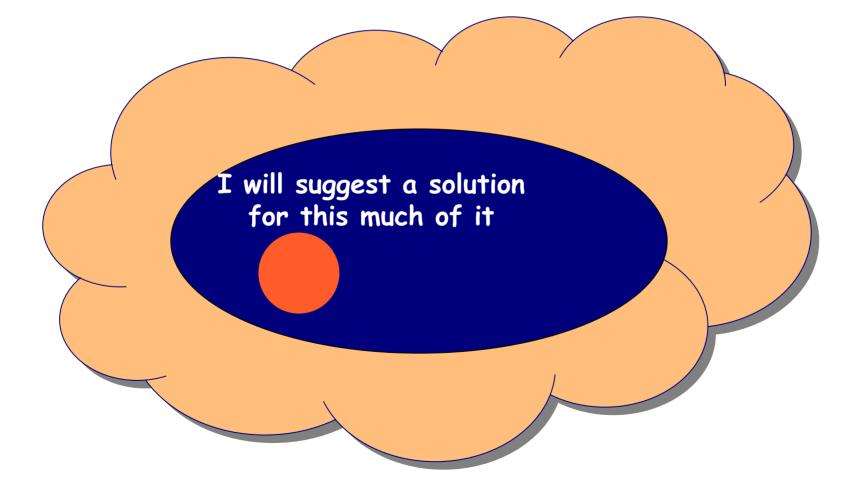
Nenad Medvidovic

Computer Science Department University of Southern California Los Angeles, U.S.A. <u>neno@usc.edu</u>

- Identify some challenges
- Suggest some solutions
- Motivate future research
- Invite dissenting opinions

The problem is this big and sometimes ill-defined





Gr

But, hopefully, we will have this much fun!

What Is Dependability?

- Degree of user confidence that the system will operate as expected
- Key dimensions
 - Availability
 - Reliability
 - Security
 - 🗆 Safety
- But also
 - Repairability
 - Maintainability
 - Survivability
 - Fault tolerance

What Is Architecture?

- A high-level model of a system
 The system's "blueprint"
- Represents system organization
 - 🗆 Data
 - Computation
 - Interaction
 - Structure
- Embodies system properties
 - Communication integrity, performance, throughput, liveness, ...
 - > Can/does it embody dependability?
 - > (how) Can those properties be transferred to the system itself?

A "Traditional" Architectural Model

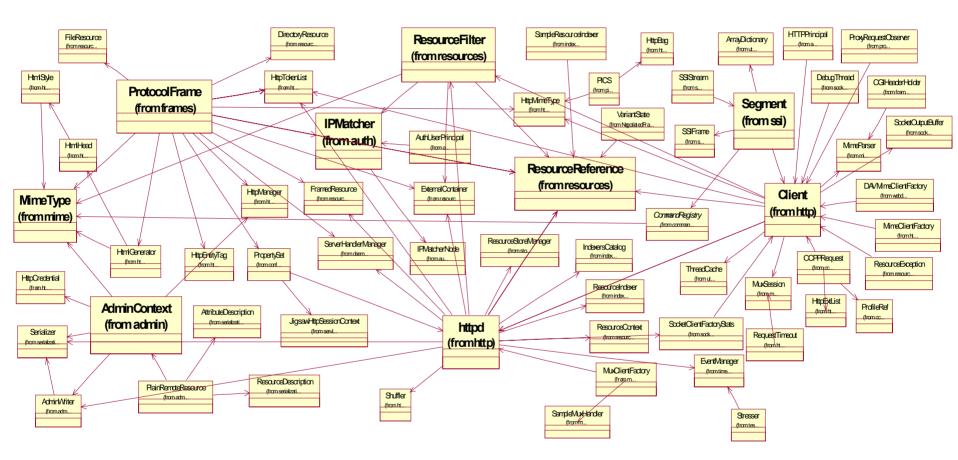
```
connector Pipe =
 role Writer = write \rightarrow Writer \square close \rightarrow \checkmark
 role Reader =
   let ExitOnly = close \rightarrow \sqrt{}
   in let DoRead = (read \rightarrow Reader
               \square read-eof \rightarrow ExitOnly)
   in DoRead Π ExitOnly
 glue = let ReadOnly = Reader.read \rightarrow ReadOnly
               [] Reader.read-eof \rightarrow Reader.close \rightarrow \checkmark
               \blacksquare Reader.close \rightarrow \sqrt{}
      in let WriteOnly = Writer.write \rightarrow WriteOnly
               🛛 Writer.close \rightarrow 🎝
      in Writer.write → glue
        ■ Reader.read → glue
        \blacksquare Writer.close \rightarrow ReadOnly
        Reader.close \rightarrow WriteOnly
```

A "Traditional" Architectural Model

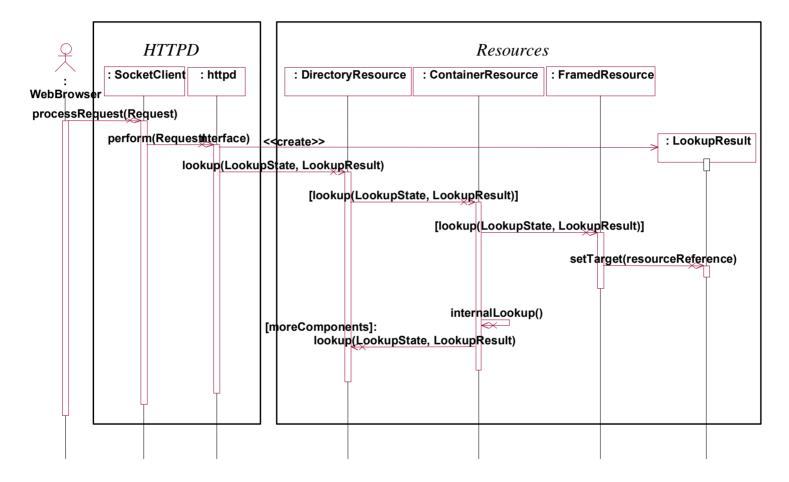
```
connector Pipe =
 role Writer = write \rightarrow Writer \square close \rightarrow \checkmark
 role Reader =
   let ExitOnly = close \rightarrow \sqrt{}
   in let DoRead = (read \rightarrow Reader
               \square read-eof \rightarrow ExitOnly)
   in DoRead Π ExitOnly
 glue = let ReadOnly = Reader.read \rightarrow ReadOnly
               [] Reader.read-eof \rightarrow Reader.close \rightarrow \checkmark
               \blacksquare Reader.close \rightarrow \sqrt{}
      in let WriteOnly = Writer.write \rightarrow WriteOnly
               🛛 Writer.close \rightarrow 🎝
      in Writer.write → glue
        ■ Reader.read → glue
        \blacksquare Writer.close \rightarrow ReadOnly
        Reader.close \rightarrow WriteOnly
```

What/where is the dependability?

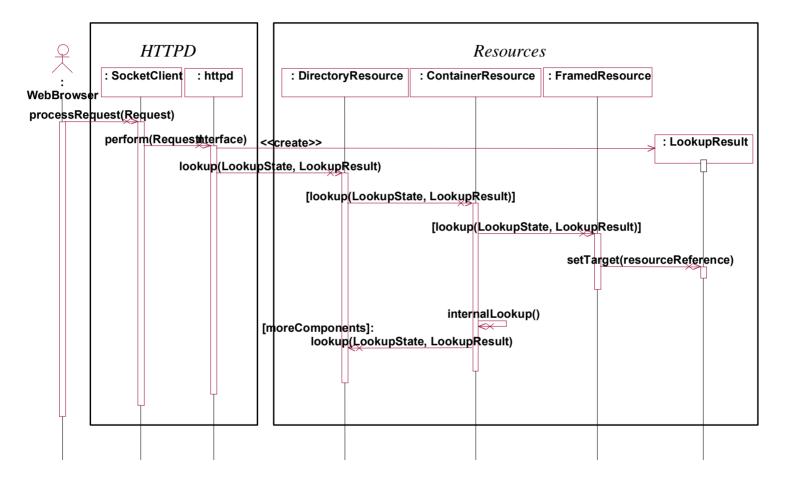
A "Standard" Architectural Model



A "Standard" Architectural Model



A "Standard" Architectural Model



What/where is the dependability?

But, We Can Model Anything

- Meta-H, ROOM, UniCon, etc. can help ensure real-time properties in models
- Markov chains can help ensure reliability in models
- Multi-versioning connectors can help ensure fault tolerance

. . .

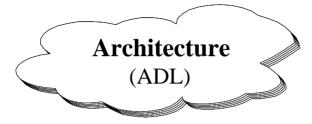
 Code/data mobility and replication formalisms can help ensure availability

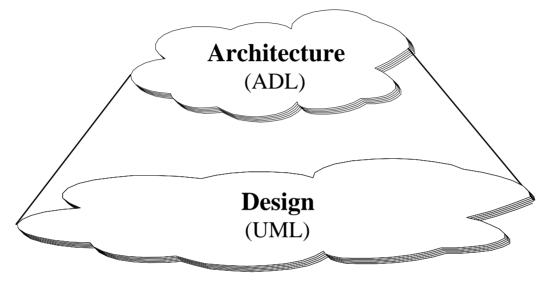
But, We Can Model Anything

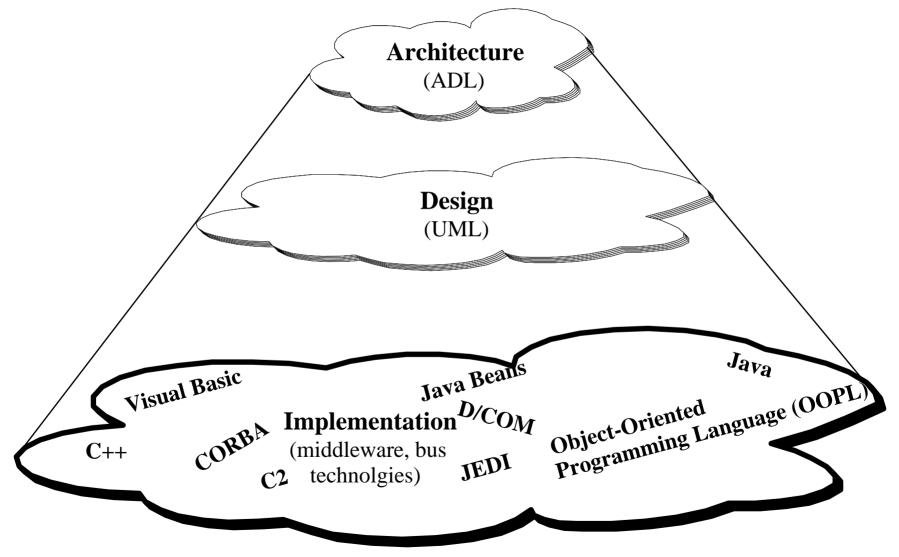
- Meta-H, ROOM, UniCon, etc. can help ensure real-time properties in models
- Markov chains can help ensure reliability in models
- Multi-versioning connectors can help ensure fault tolerance
- Code/data mobility and replication formalisms can help ensure availability

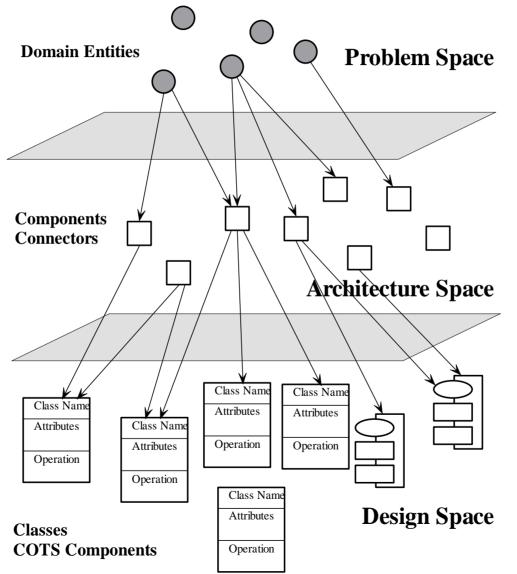
So then the problem is solved, right?

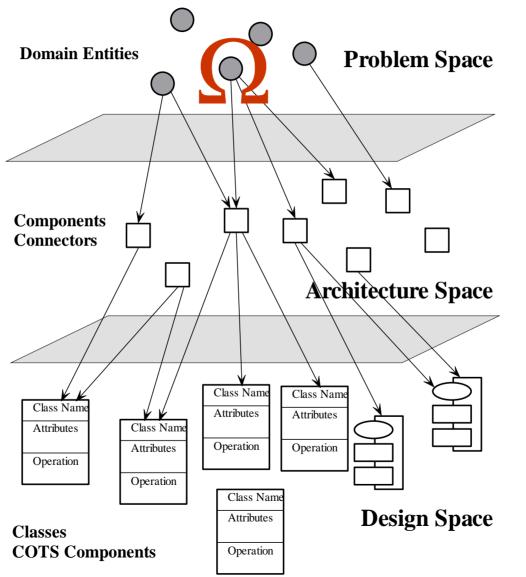
. . .

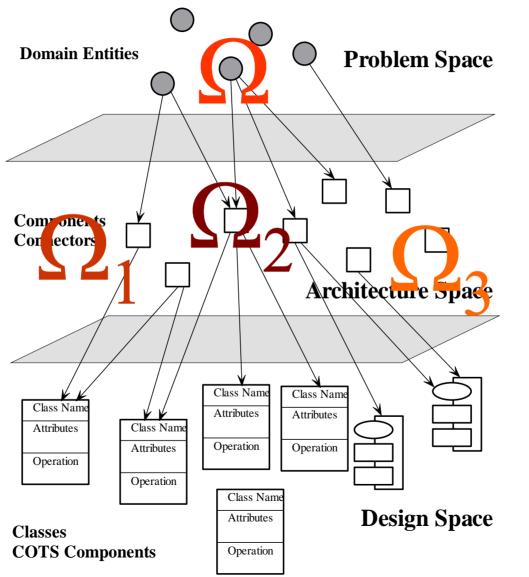


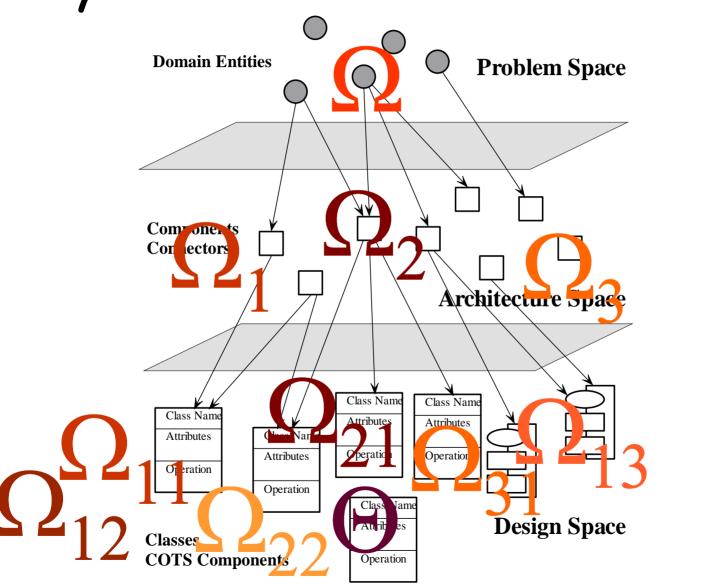




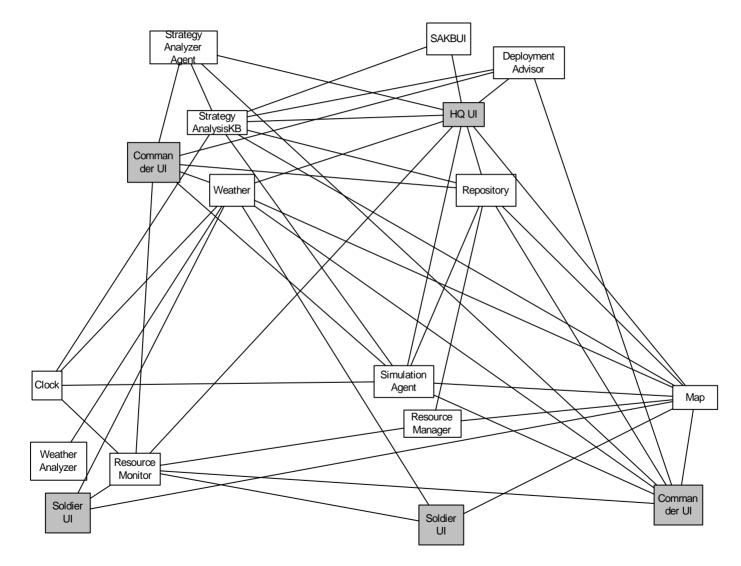




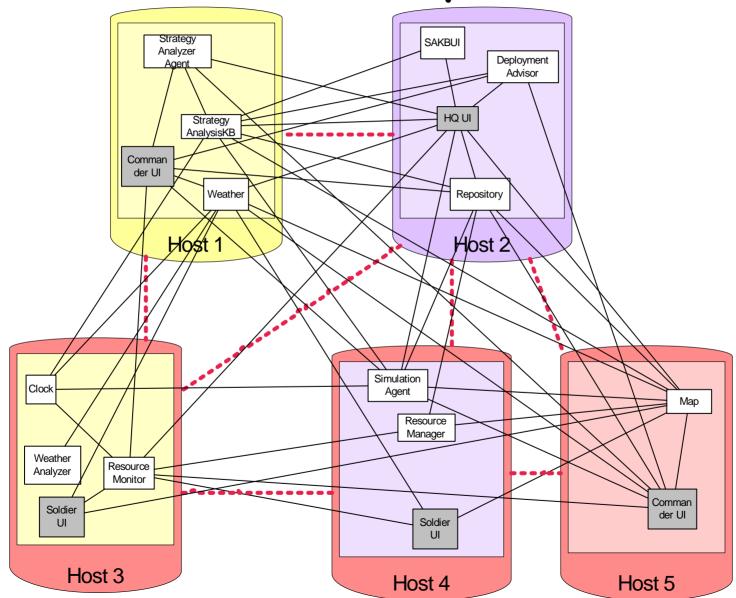




From Models to Systems



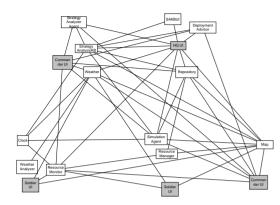
From Models to Systems



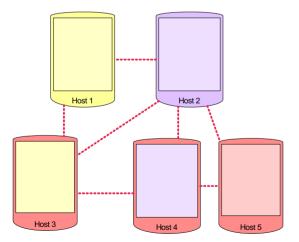
From Models to Systems

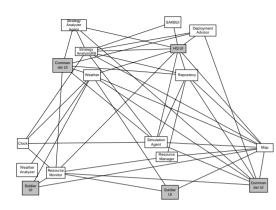


The remainder of this talk will focus on two key questions:

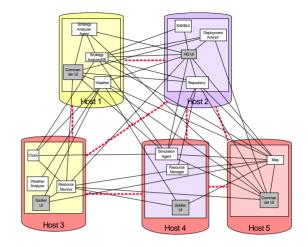


and

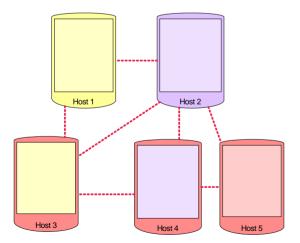


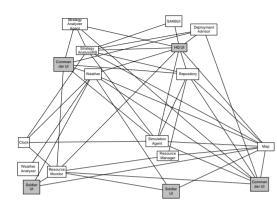


to

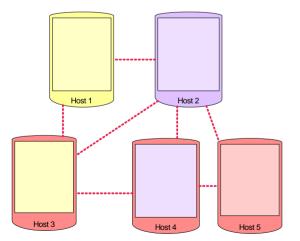


and

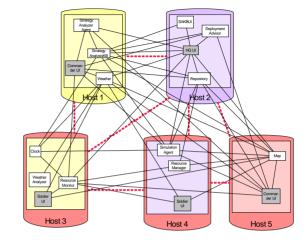




and





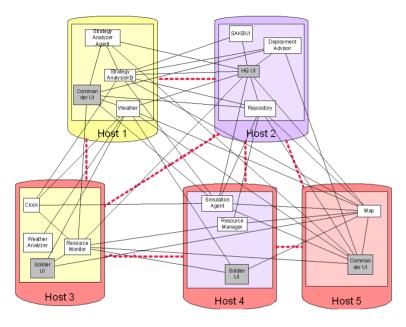


to

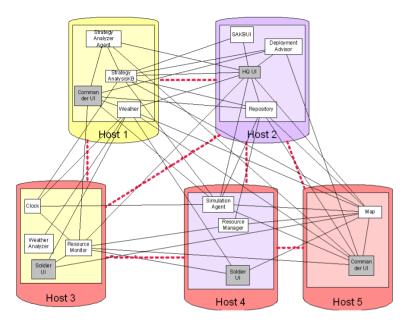


2. How do we know

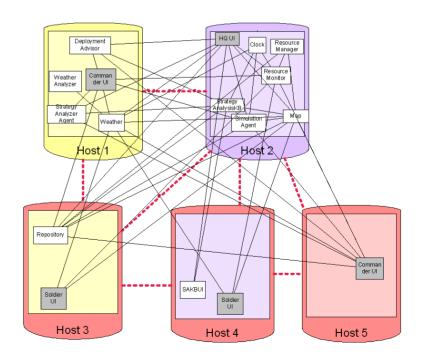
```
2. How do we know
```



```
2. How do we know
```



is "better" than



Outline

From architectures to systems
 Ensuring dependability

 Problem definition
 Proposed solution

 Concluding remarks

Outline

From architectures to systems
 Ensuring dependability

 Problem definition
 Proposed solution

 Concluding remarks

How Do I Dependably Implement an Architecture?

- Architectures provide *high-level* concepts
 - Components, connectors, ports, events, configurations
- Programming languages provide *low-level* constructs
 Variables, arrays, pointers, procedures, objects
- Bridging the two often is an art-form
 - □ Middleware can help "split the difference"
- Existing middleware technologies
 - Support some architectural concepts (e.g., components, events)
 - but not others (e.g., connectors, configurations)
 - Impose particular architectural styles

How Do I Dependably Implement an Architecture?

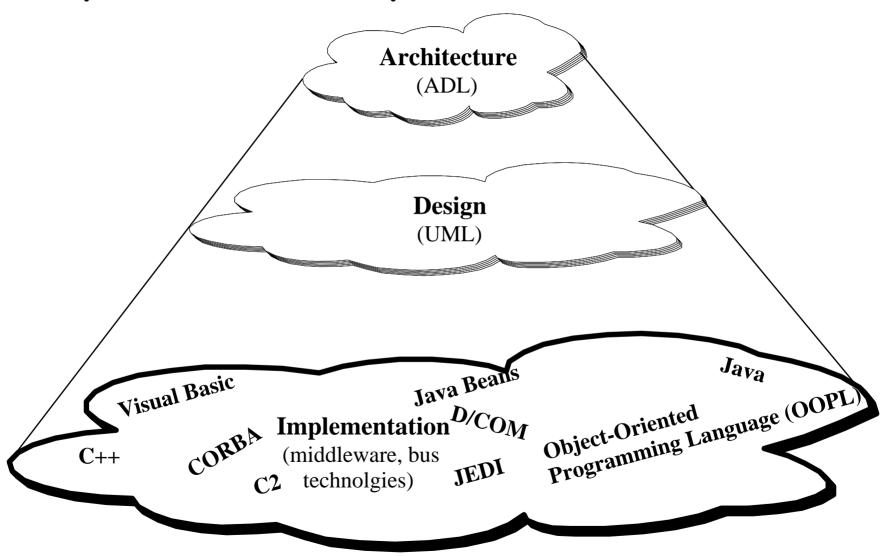
- Architectures provide *high-level* concepts
 - Components, connectors, ports, events, configurations
- Programming languages provide *low-level* constructs
 Variables, arrays, pointers, procedures, objects
- Bridging the two often is an art-form
 - □ Middleware can help "split the difference"
- Existing middleware technologies
 - Support some architectural concepts (e.g., components, events)
 - but not others (e.g., connectors, configurations)
 - Impose particular architectural styles

What is needed is "architectural middleware"

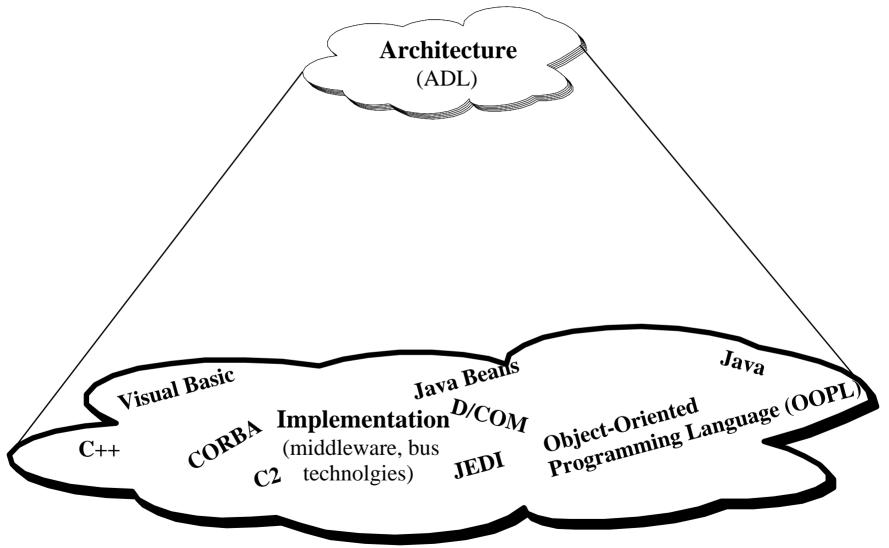
Architectural Middleware

- Natively support architectural concepts as middleware constructs
- Include system design support
 - $\hfill\square$ Typically via an accompanying ADL and analysis tools
 - May support explicit architectural styles
- Support round-trip development
 - $\hfill\square$ From architecture to implementation and back
- Support automated transformation of architectural models to implementations
 - □ i.e., dependable implementation
- Examples
 - 🗆 ArchJava
 - 🗆 Aura
 - 🗆 c2.fw
 - Prism-MW

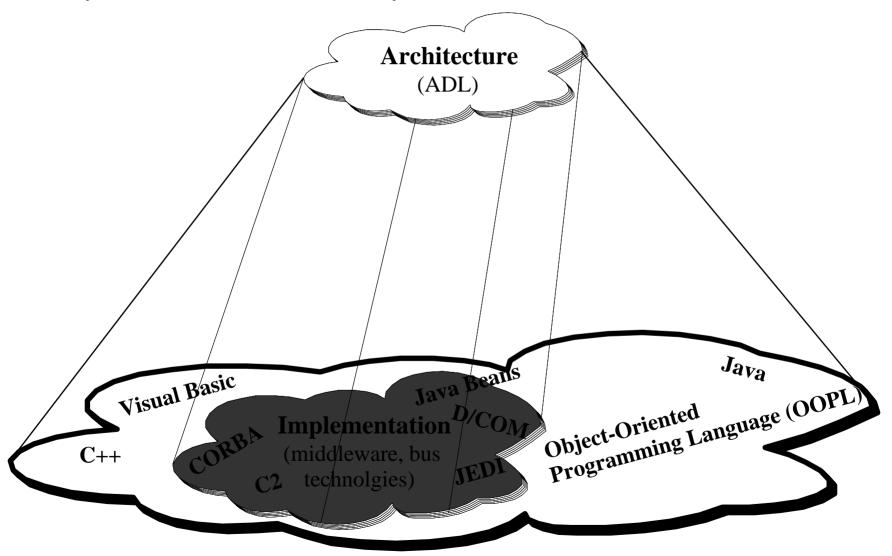
Dependable Implementation

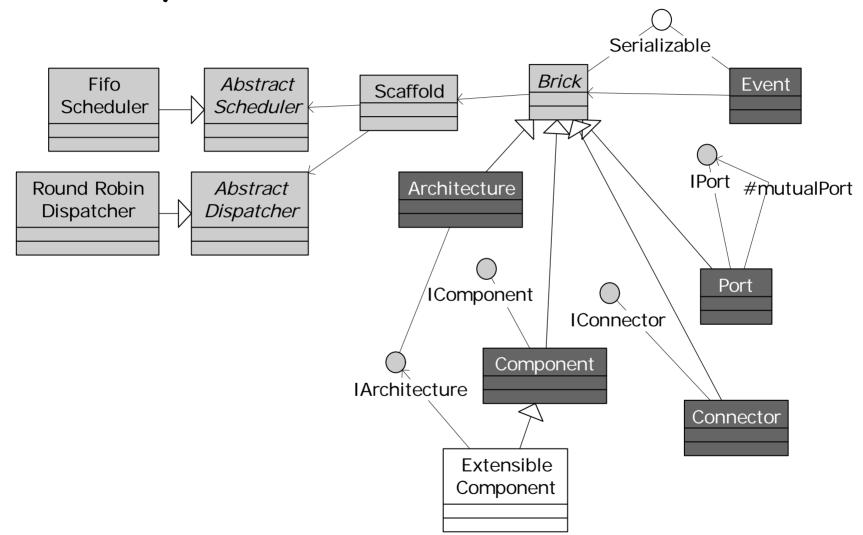


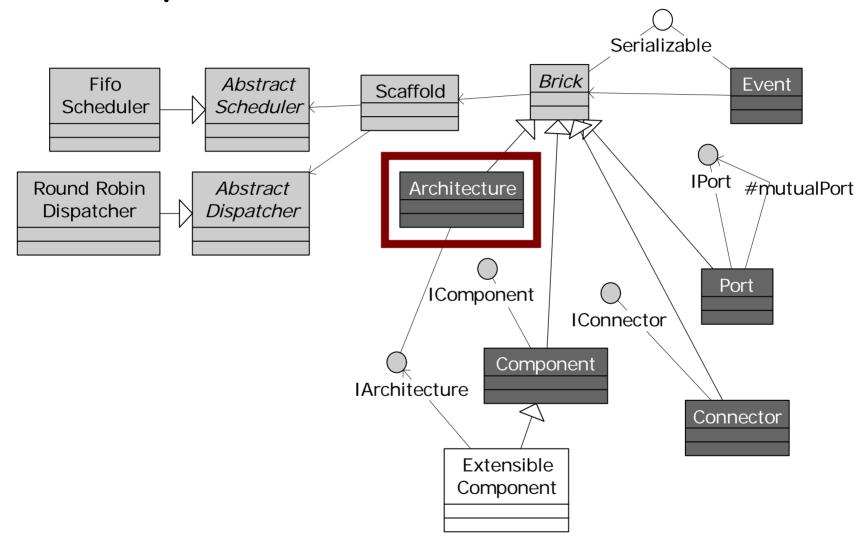
Dependable Implementation

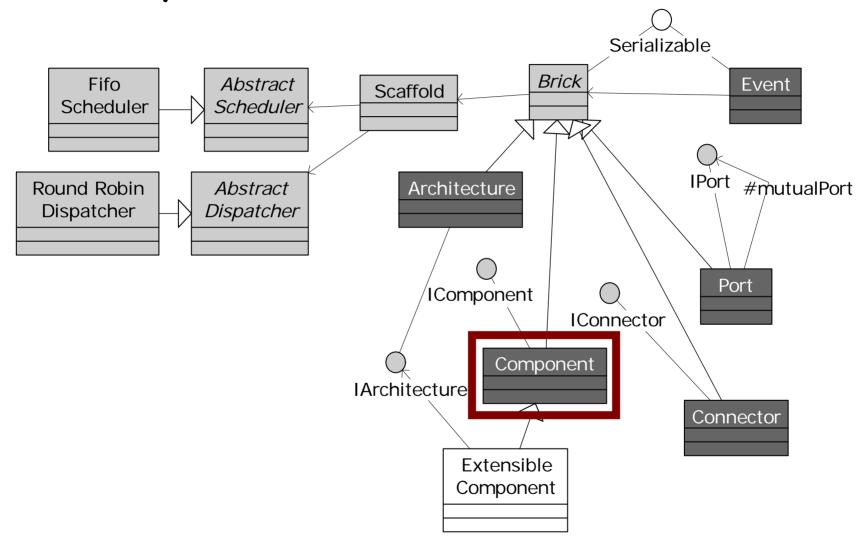


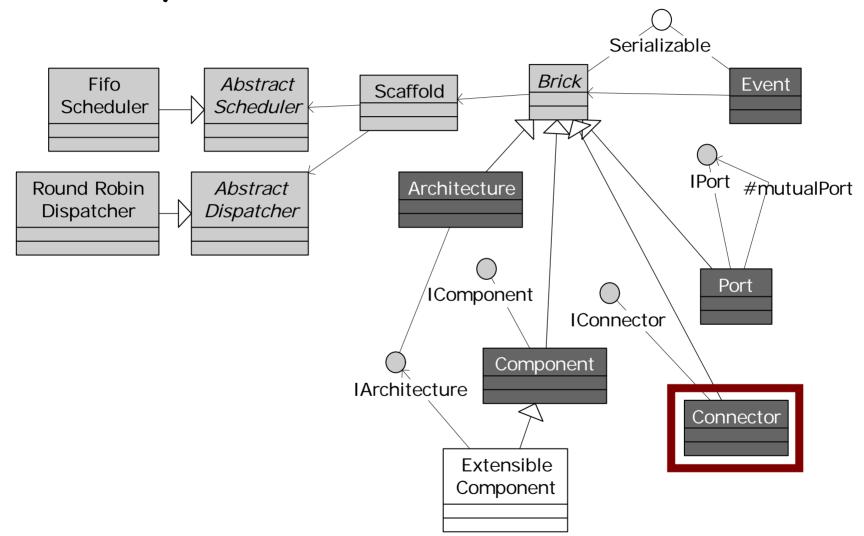
Dependable Implementation

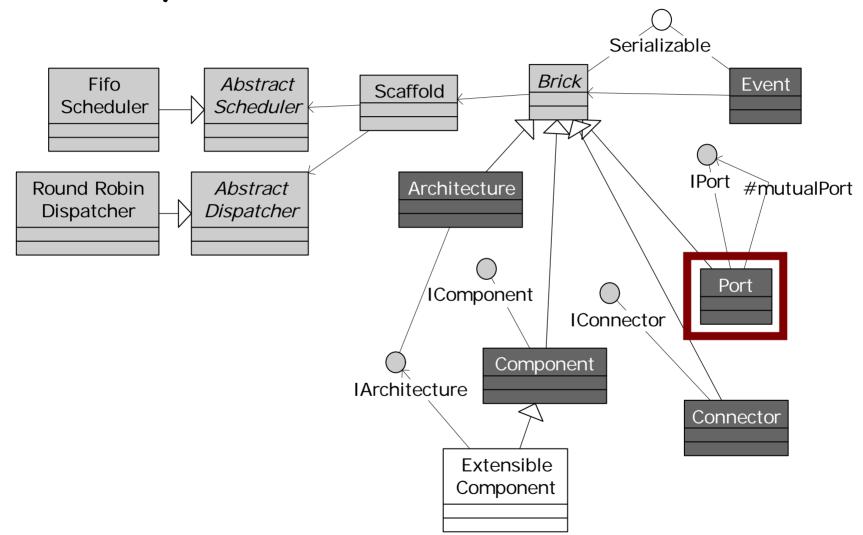


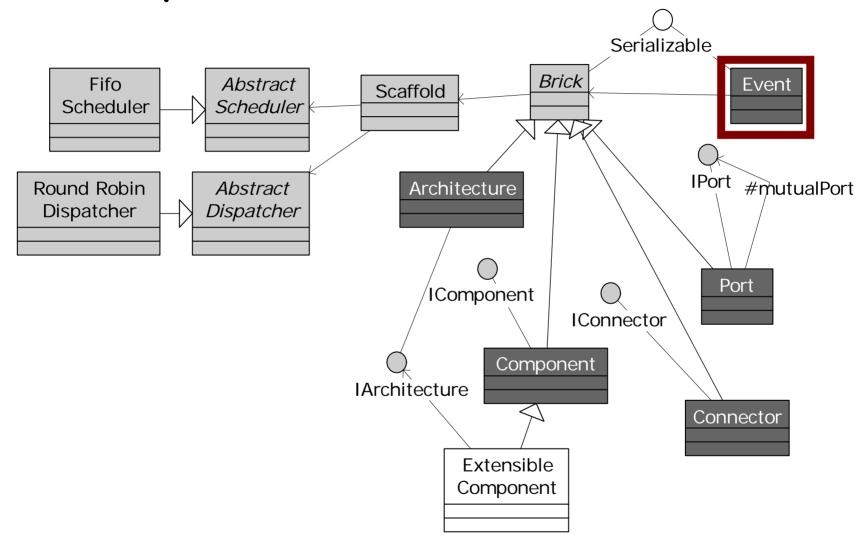


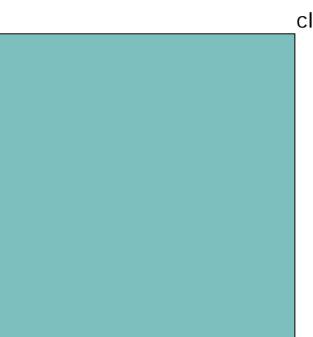






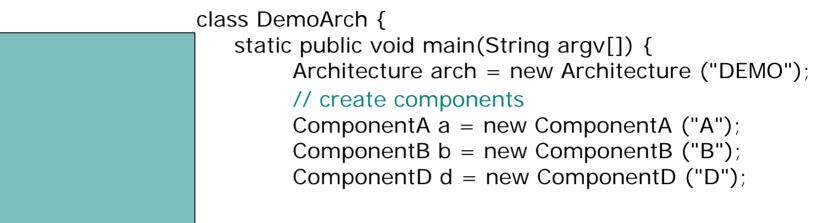






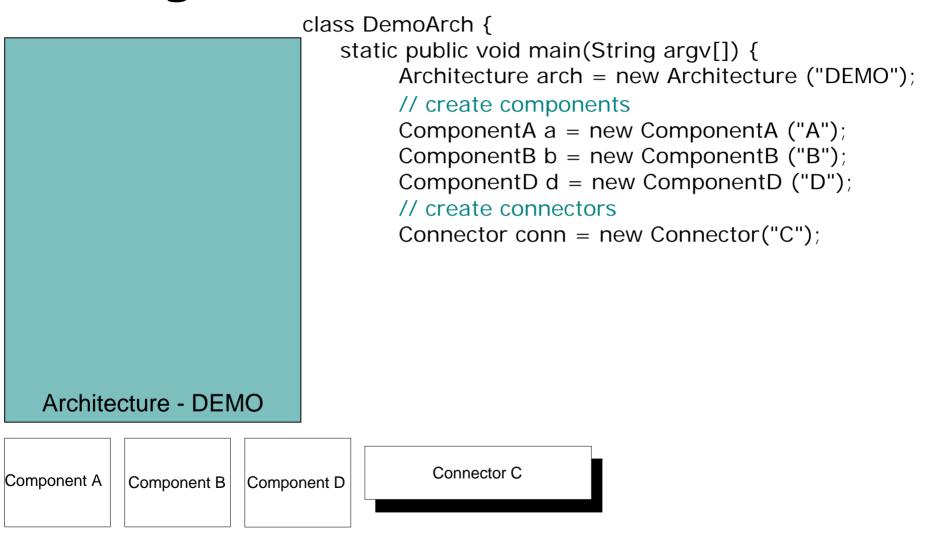
class DemoArch {
 static public void main(String argv[]) {
 Architecture arch = new Architecture ("DEMO");
 }
}

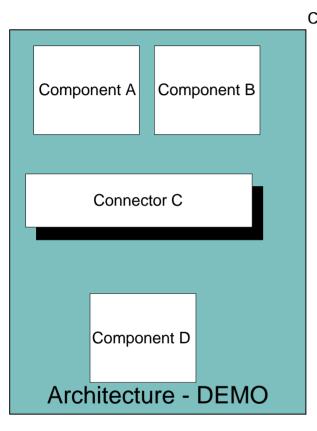
Architecture - DEMO



Architecture - DEMO

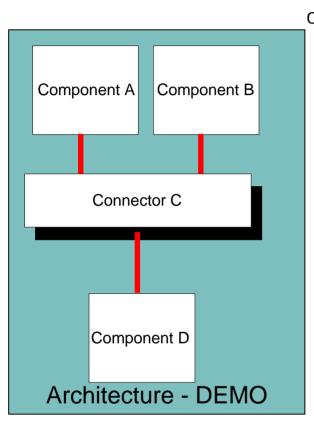
Component A Component B Component D





class DemoArch { static public void main(String argv[]) { Architecture arch = new Architecture ("DEMO"); // create components ComponentA a = new ComponentA ("A"); ComponentB b = new ComponentB ("B"); ComponentD d = new ComponentD ("D"); // create connectors Connector conn = new Connector("C"); // add components and connectors arch.addComponent(a); arch.addComponent(b); arch.addComponent(d); arch.addConnector(conn);

}



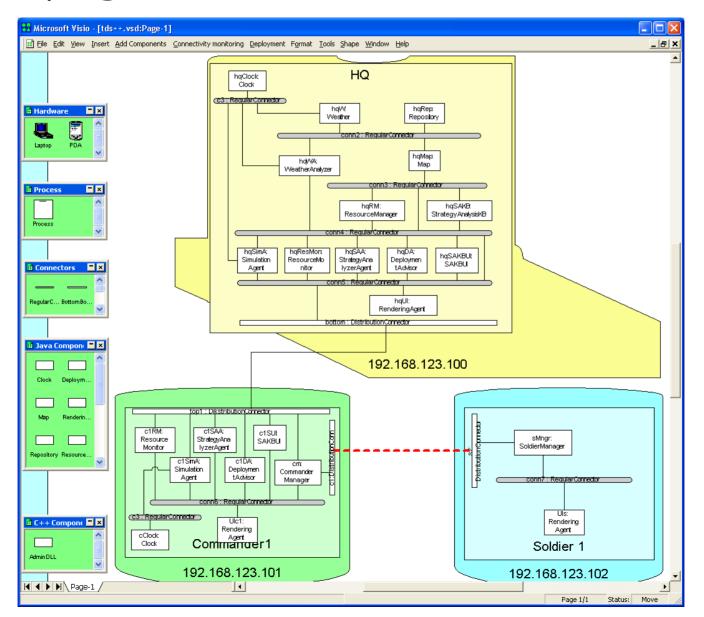
class DemoArch {
 static public void main(String argv[]) {
 Architecture arch = new Architecture ("DEMO");
 // create components
 ComponentA a = new ComponentA ("A");
 ComponentB b = new ComponentB ("B");
 ComponentD d = new ComponentD ("D");
 // create connectors
 Connector conn = new Connector("C");
 // add components and connectors
 }
}

arch.addComponent(a); arch.addComponent(b); arch.addComponent(d); arch.addConnector(conn);

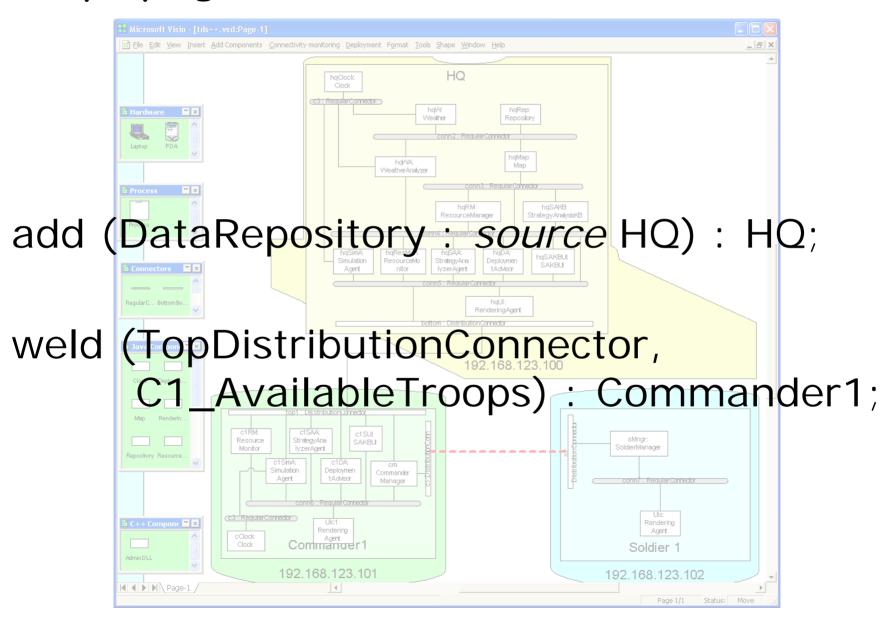
// establish the interconnections

arch.weld(a, conn); arch.weld(b, conn); arch.weld(conn, d)

Deploying a Prism-MW Architecture



Deploying a Prism-MW Architecture



Outline

From architectures to systems
 Ensuring dependability

 Problem definition
 Proposed solution

 Concluding remarks

Outline

From architectures to systems
 Ensuring dependability
 Problem definition

 Proposed solution

 Concluding remarks

Availability

- The degree to which a system is operational and accessible when required for use [IEEE]
- Deployment architecture influences availability
 - Components on the same host can communicate regardless of the network's status
 - Components on different hosts are insulated from each other's failures
- Quantifying availability
 - \Box Ratio of the

number of successfully completed interactions in the system to the

total number of attempted interactions

Maximizing Availability a priori

- We may not know many relevant system parameters
 - Dependability of each component
 - Frequency of component interactions
 - Volume of component interactions
 - Dependability of component interactions
 - □ CPU usage on each host
 - Dependability of each host
 - Effective bandwidth of each network connection
 - Dependability of each network connection

□...

Maximizing Availability a priori

- We may not know many relevant system parameters
 - Dependability of each component
 - Frequency of component interactions
 - Volume of component interactions
 - Dependability of component interactions
 - □ CPU usage on each host
 - Dependability of each host
 - Effective bandwidth of each network connection
 - Dependability of each network connection

□...

The current deployment architecture may not work well

Given:

- (1) a set C of n components (n = |C|), a relation
- $\mathit{freq}: C \times C \to \Re$, and a function $\mathit{mem}_{\mathit{comp}}: C \to \Re$

- Given:
- (1) a set *C* of *n* components (n = |C|), a relation *freq* : $C \times C \rightarrow \Re$, and a function *mem*_{comp} : $C \rightarrow \Re$

$$freq(c_i, c_j) = \begin{cases} 0 & \text{if } c_i = c_j \\ frequency & \text{of comm between } c_i & \text{and } c_j & \text{if } c_i \neq c_j \end{cases}$$

(1) a set *C* of *n* components (n = |C|), a relation *freq* : $C \times C \rightarrow \Re$, and a function $mem_{comp} : C \rightarrow \Re$

$$freq(c_i, c_j) = \begin{cases} 0 & \text{if } c_i = c_j \\ frequency & \text{of comm between } c_i & \text{and } c_j & \text{if } c_i \neq c_j \end{cases}$$

 $mem_{comp}(c) = required memory for c$

Given:

(2) a set H of k hardware nodes (k = |H|), a relation

 $rel: H \times H \to \Re$, and a function $mem_{host}: H \to \Re$

(2) a set H of k hardware nodes (k = |H|), a relation

 $rel: H \times H \rightarrow \Re$, and a function $mem_{host}: H \rightarrow \Re$

 $rel(h_i, h_j) = \begin{cases} 1 & if \quad h_i = h_j \\ 0 & if \quad h_i \text{ is not connected to } h_j \\ reliability \text{ of the link between } h_i \text{ and } h_j \text{ if } h_i \neq h_j \end{cases}$

(2) a set H of k hardware nodes (k = |H|), a relation

 $rel: H \times H \to \Re$, and a function $mem_{host}: H \to \Re$

$$rel(h_i, h_j) = \begin{cases} 1 & if \quad h_i = h_j \\ 0 & if \quad h_i \text{ is not connected to } h_j \\ reliability \text{ of the link between } h_i \text{ and } h_j \text{ if } h_i \neq h_j \end{cases}$$

 $mem_{host}(h) = available memory on host h$

Given:

- (3) Two relations that restrict locations of software components
 - $loc: C \times H \rightarrow \{0,1\}$ $colloc: C \times C \rightarrow \{-1,0,1\}$

Given:

(3) Two relations that restrict locations of software components $loc: C \times H \rightarrow \{0,1\}$ $colloc: C \times C \rightarrow \{-1,0,1\}$

 $loc(c_i, h_j) = \begin{cases} 1 & if \quad c_i \text{ can be deployed onto } h_j \\ 0 & if \quad c_i \text{ cannot be deployed onto } h_j \end{cases}$

Given:

(3) Two relations that restrict locations of software components

 $loc: C \times H \to \{0,1\} \qquad colloc: C \times C \to \{-1,0,1\}$

 $loc(c_i, h_j) = \begin{cases} 1 & if \quad c_i \text{ can be deployed onto } h_j \\ 0 & if \quad c_i \text{ cannot be deployed onto } h_j \end{cases}$

 $colloc(c_i, c_j) = \begin{cases} -1 & if \quad c_i \text{ cannot be on the same host as } c_j \\ 1 & if \quad c_i \text{ has to be on the same host as } c_j \\ 0 & if \quad there \text{ are no restrictions on collocation of } c_i \text{ and } c_j \end{cases}$

Find a function $f: C \rightarrow H$ such that the

system's overall availability

$$A = \frac{\sum_{i=1}^{n} \sum_{j=1}^{n} \left(freq(c_{i}, c_{j}) * rel(f(c_{i}), f(c_{j})) \right)}{\sum_{i=1}^{n} \sum_{j=1}^{n} freq(c_{i}, c_{j})}$$

is maximized

Find a function $f: C \rightarrow H$ such that the

system's overall availability

$$A = \frac{\sum_{i=1}^{n} \sum_{j=1}^{n} \left(freq(c_{i}, c_{j}) * rel(f(c_{i}), f(c_{j})) \right)}{\sum_{i=1}^{n} \sum_{j=1}^{n} freq(c_{i}, c_{j})}$$

is maximized, and the following three conditions are satisfied:

Find a function $f: C \rightarrow H$ such that the

system's overall availability

$$A = \frac{\sum_{i=1}^{n} \sum_{j=1}^{n} \left(freq(c_{i}, c_{j}) * rel(f(c_{i}), f(c_{j})) \right)}{\sum_{i=1}^{n} \sum_{j=1}^{n} freq(c_{i}, c_{j})}$$

is maximized, and the following three conditions are satisfied:

$$\forall i \in [1,k] \left\{ \forall j \in [1,n] \quad f(c_j) = h_i \quad \left| \begin{array}{c} \sum_{j} mem_{comp}(c_j) \right| \leq mem_{host}(h_i) \right\}$$

Find a function $f: C \rightarrow H$ such that the

system's overall availability

$$A = \frac{\sum_{i=1}^{n} \sum_{j=1}^{n} \left(freq(c_{i}, c_{j}) * rel(f(c_{i}), f(c_{j})) \right)}{\sum_{i=1}^{n} \sum_{j=1}^{n} freq(c_{i}, c_{j})}$$

is maximized, and the following three conditions are satisfied:

$$\forall j \in [1, n] \qquad loc(c_j, f(c_j)) = 1$$

Find a function $f: C \rightarrow H$ such that the

system's overall availability

$$A = \frac{\sum_{i=1}^{n} \sum_{j=1}^{n} \left(freq(c_{i}, c_{j}) * rel(f(c_{i}), f(c_{j})) \right)}{\sum_{i=1}^{n} \sum_{j=1}^{n} freq(c_{i}, c_{j})}$$

is maximized, and the following three conditions are satisfied:

 $\forall k \in [1, n] \quad \forall l \in [1, n]$ $(colloc (c_k, c_l) = 1) \Rightarrow (f(c_k) = f(c_l))$ $(colloc (c_k, c_l) = -1) \Rightarrow (f(c_k) \neq f(c_l))$

Find a function $f: C \rightarrow H$ such that the

system's overall availability

$$A = \frac{\sum_{i=1}^{n} \sum_{j=1}^{n} \left(freq(c_{i}, c_{j}) * rel(f(c_{i}), f(c_{j})) \right)}{\sum_{i=1}^{n} \sum_{j=1}^{n} freq(c_{i}, c_{j})}$$

is maximized, and the following three conditions are satisfied:

Note that the possible number of different functions f is k^n

Outline

From architectures to systems
 Ensuring dependability

 Problem definition
 Proposed solution

 Concluding remarks

Overview of the Approach

Objective

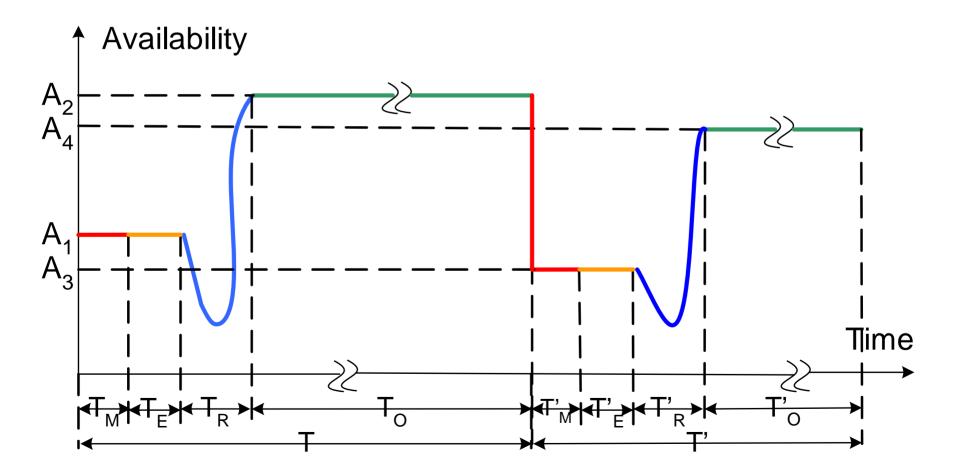
- Identify the problem
 - Log and examine system events
 - Actively monitor the system during runtime
- Develop a solution
 - Decide which data to cache
 - Decide which components to replicate
 - Introduce multiple execution modes
 - Calculate an improved system deployment
- Apply the solution to eliminate the problem
 - Cache or hoard data
 - Replicate data or code
 - Redeploy (parts of) the system

Overview of the Approach

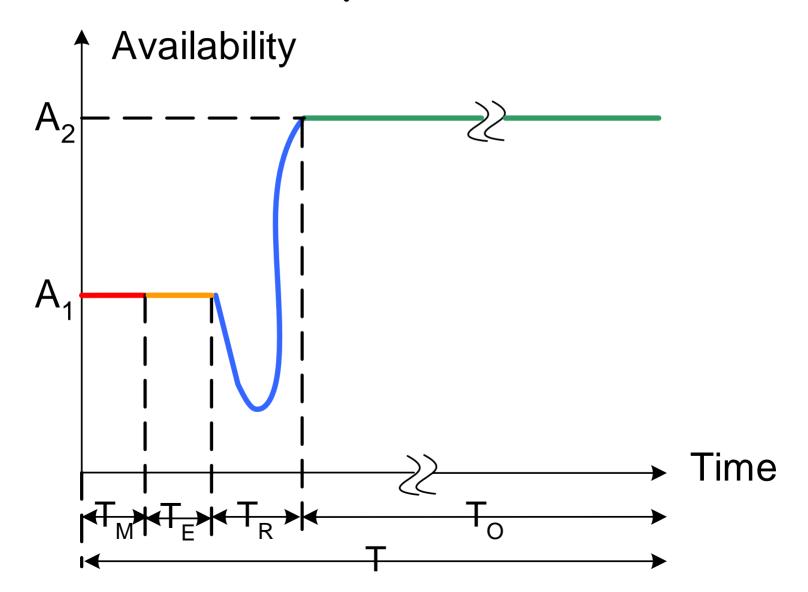
Objective

- Identify the problem
 - Log and examine system events
 - > Actively monitor the system during runtime
- Develop a solution
 - Decide which data to cache
 - Decide which components to replicate
 - Introduce multiple execution modes
 - > Calculate an improved system deployment
- Apply the solution to eliminate the problem
 - Cache or hoard data
 - Replicate data or code
 - Redeploy (parts of) the system

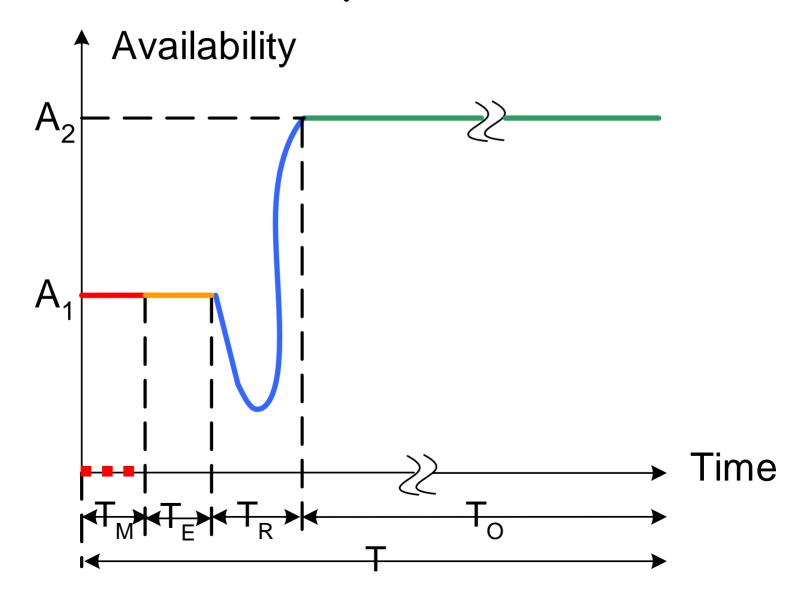
Improving Availability via Redeployment



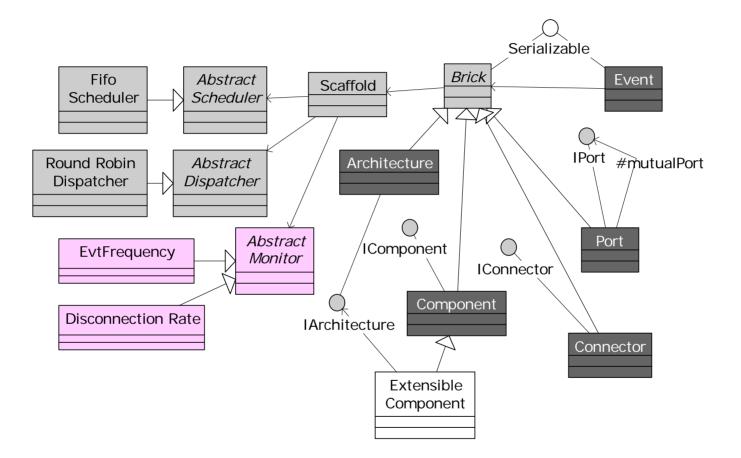
First Identify the Problem



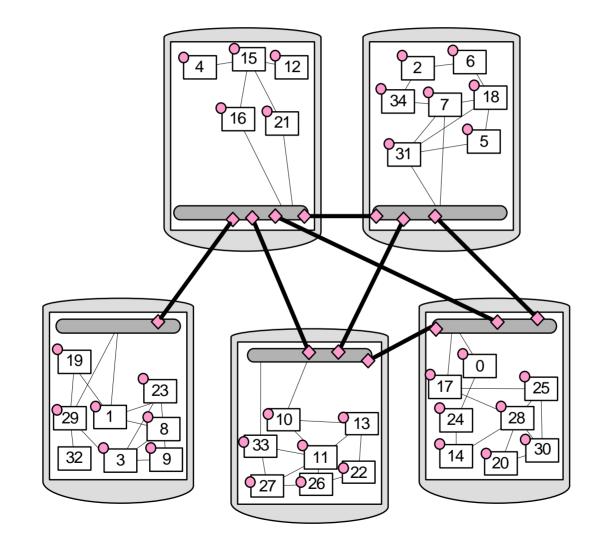
First Identify the Problem

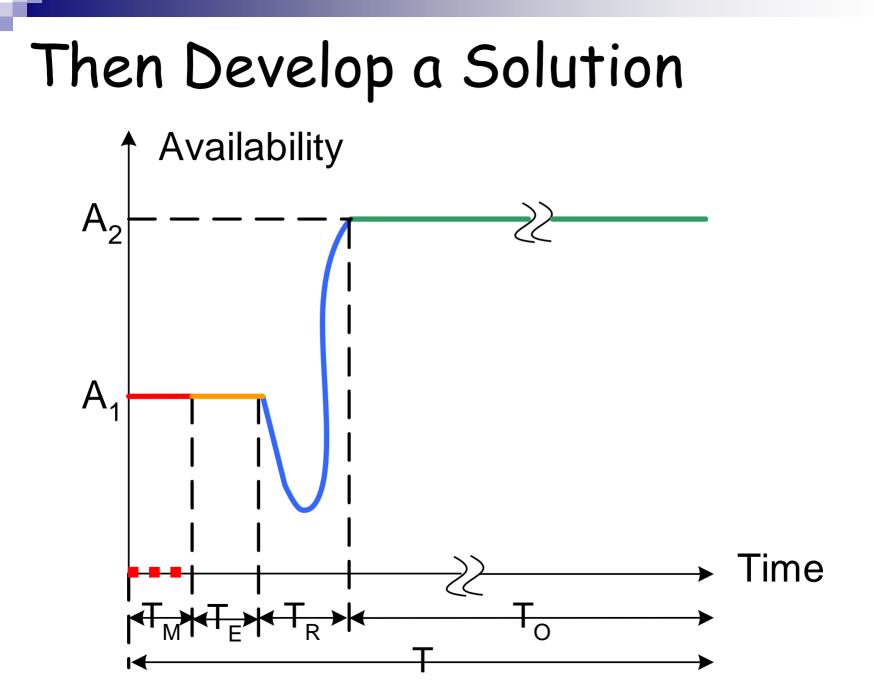


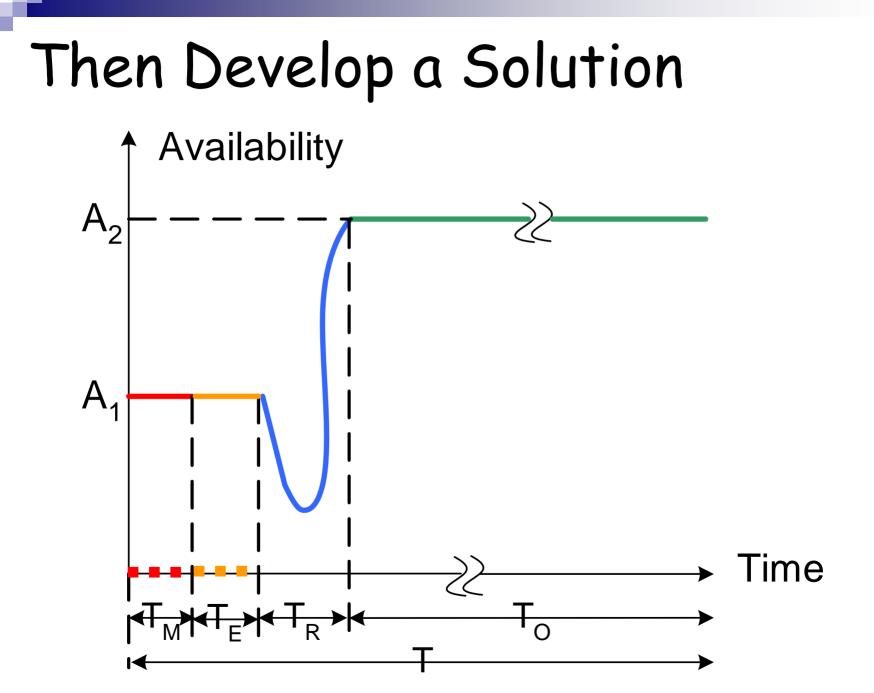
Monitoring in Prism-MW



Monitoring in Action







Estimating in Action

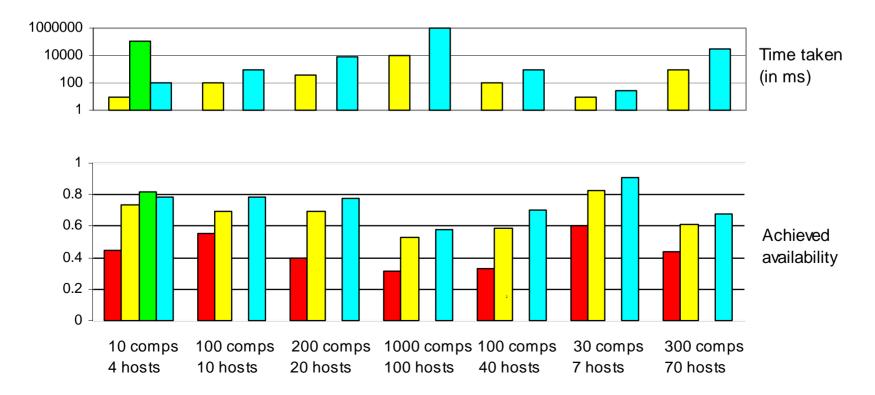
Suite of algorithms:

Exact - exponential complexity

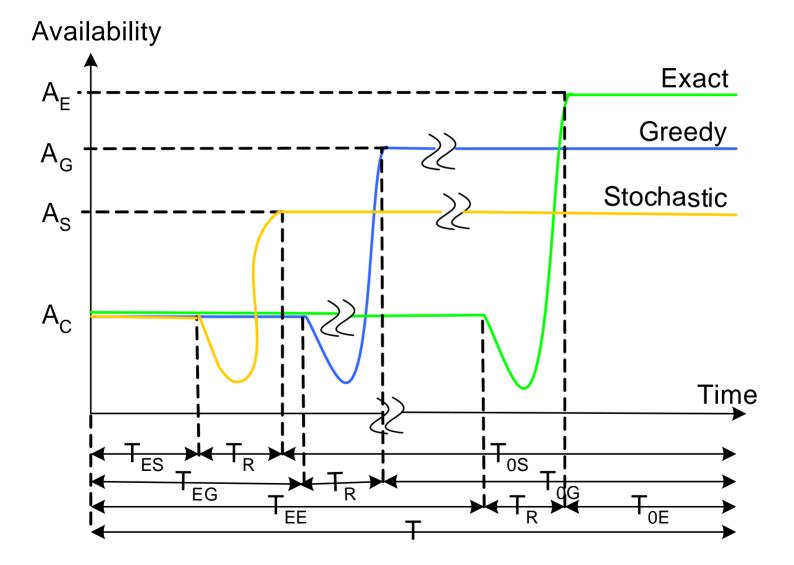
Stochastic – quadratic complexity

Adaptive greedy – cubic complexity

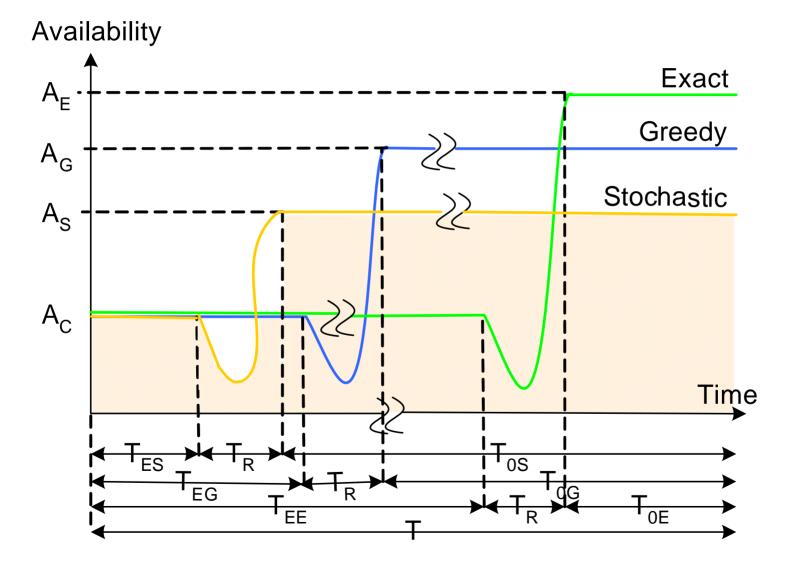
Decentralized - quadratic complexity

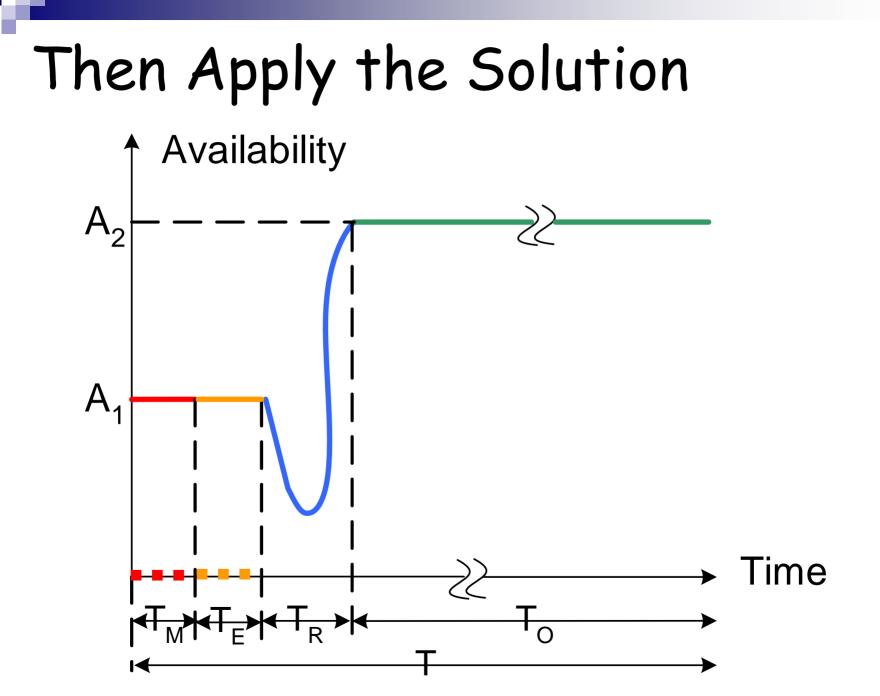


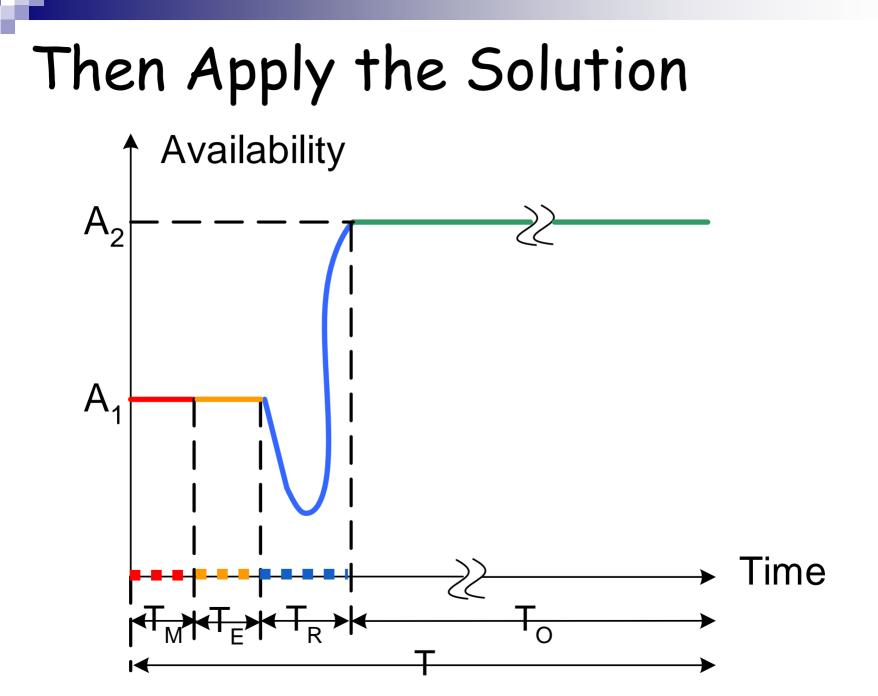
Automatic Algorithm Selection



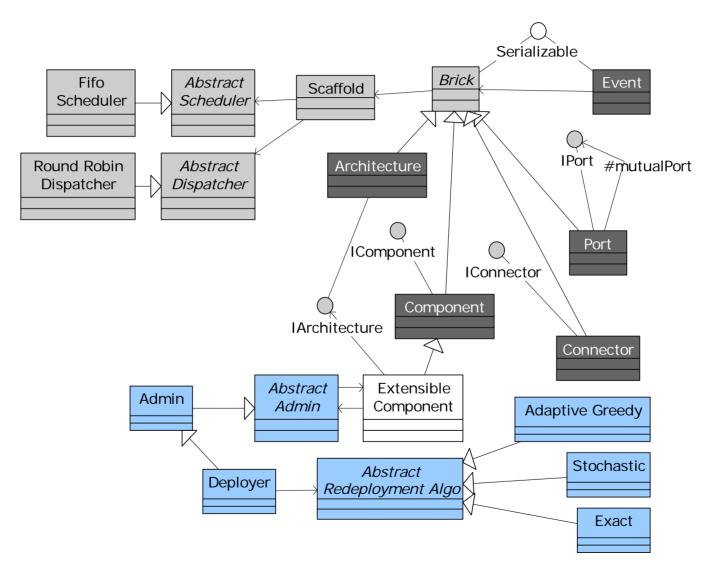
Automatic Algorithm Selection



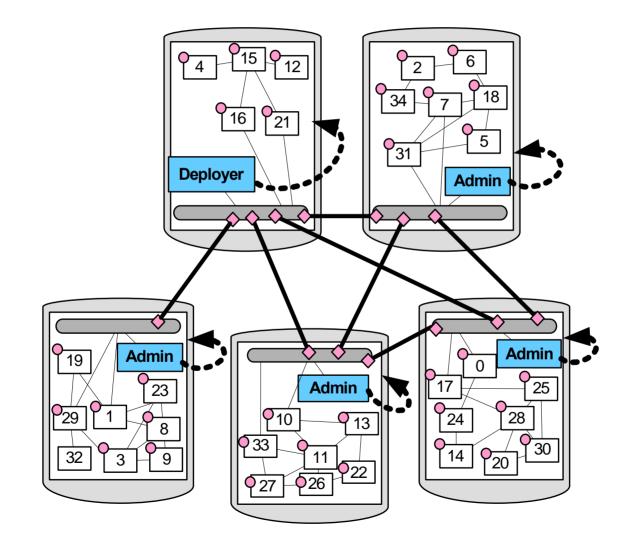


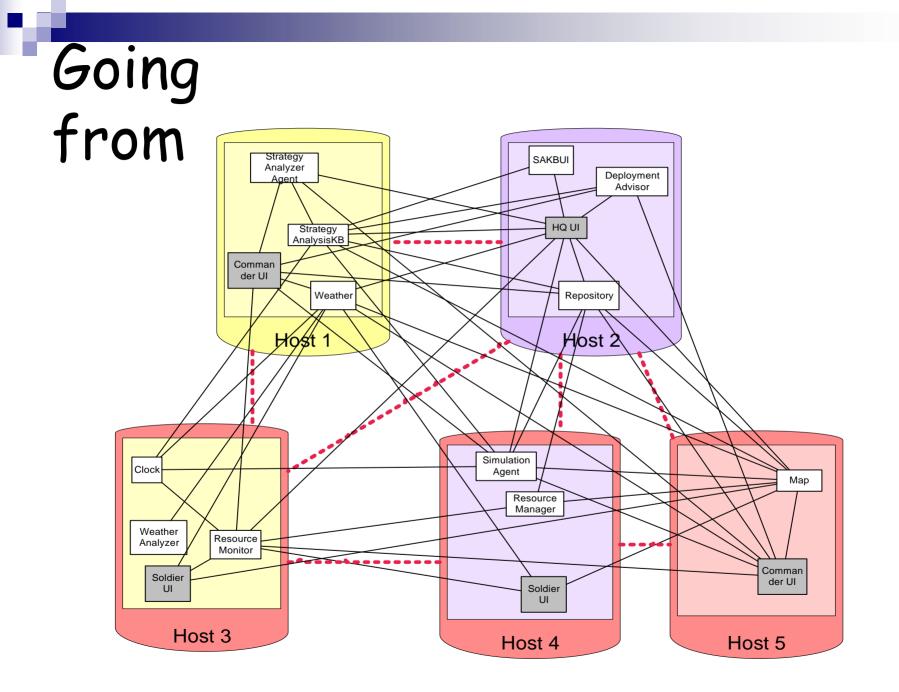


Redeployment in Prism-MW

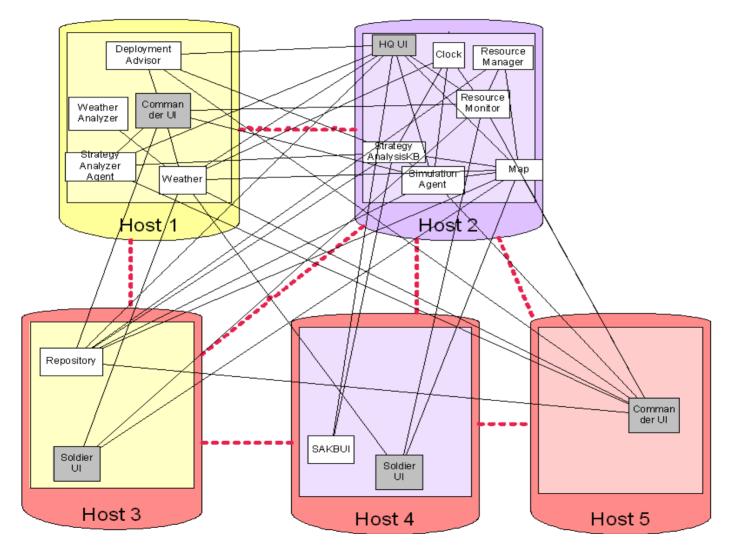


Redeployment in Action





to



Outline

From architectures to systems
 Ensuring dependability

 Problem definition
 Proposed solution

 Concluding remarks

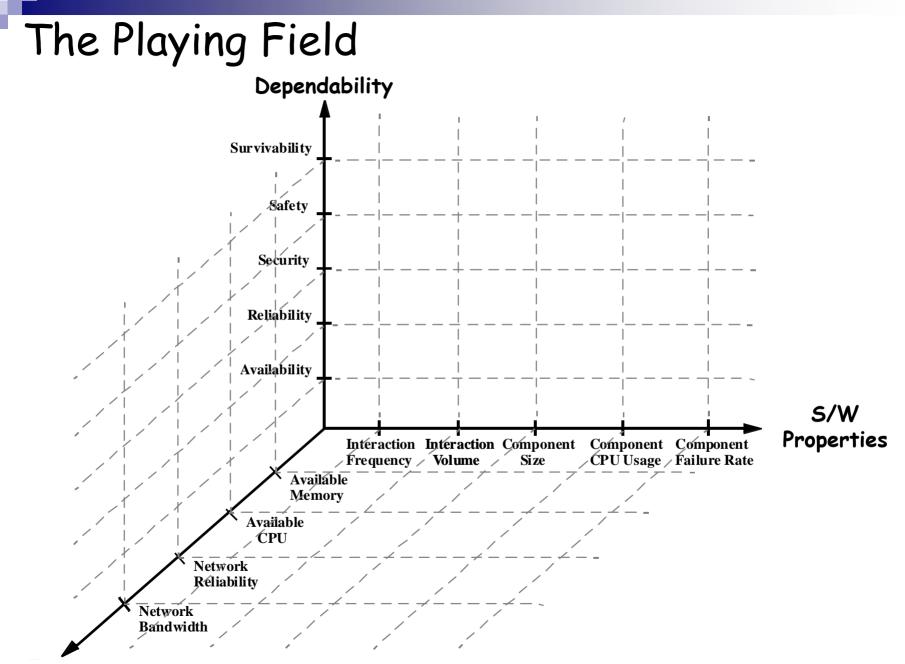
Concluding Remarks

Still so much to do

- Enriching/completing the models
- Focusing on additional aspects of dependability
- Addressing feature interactions
- Addressing emergent properties
- Determining which concerns are (not) architectural

Promise or Illusion?

- □ If we can define it, we should be able to
 - Analyze for it
 - Build it
 - Measure it
 - Act on it
- □ So, why haven't we done it yet?



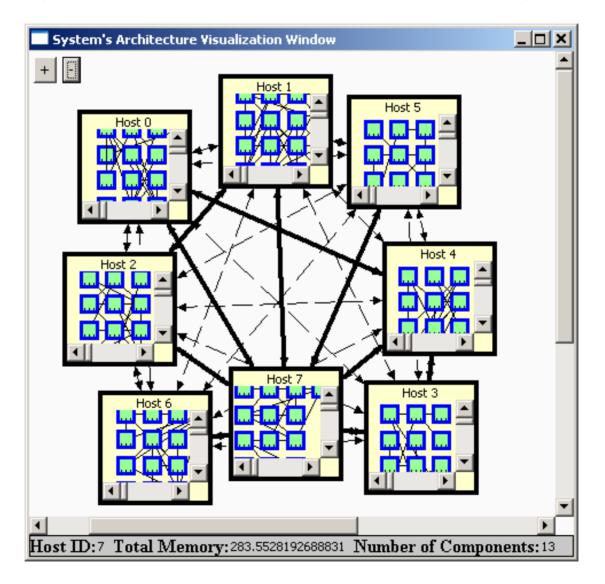
H/W Properties

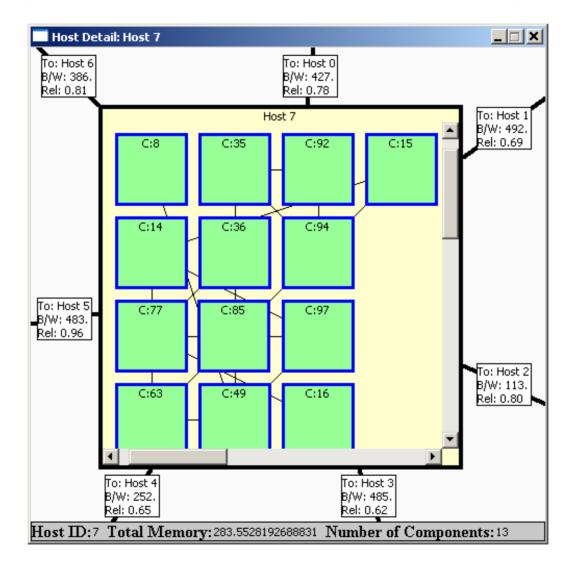
Deployment Control Window

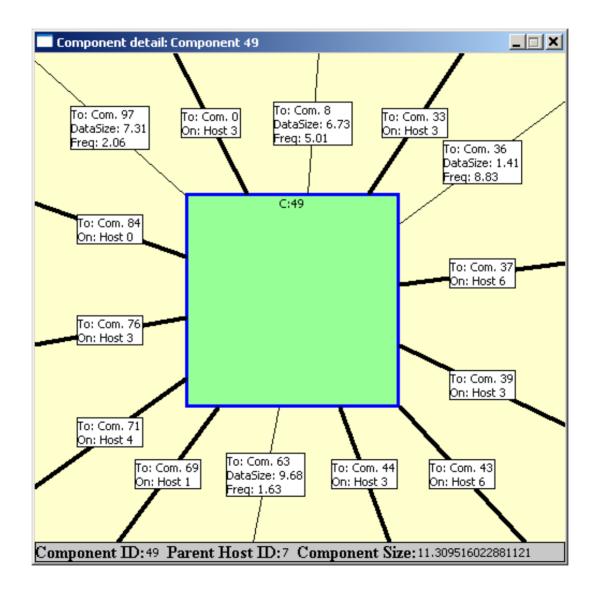
Input		Constraints					-Algor	rithms				
Number of components:	100	Components	Hosts				What	: do yo	ou want to do	o: Just	run	-
Number of hosts:	8	Component-0	Host-	0	On the same host					lust	run	
Minimum comp. memory (in KB):	10	Component-1	Host-			-11				Exa Run a	and preview	
Maximum comp. memory (in KB):	20	Component-2 Component-3	Host- Host-		Not on the same hos	;t			Un	Run a	and effect astic	
Minimum host memory (in KB):	200	Component-4	Host-		Fix to host							
Maximum host memory (in KB):	300	Component-5	Host-		- Tix to host				Bi	iased Stocha:	stic	
Minimum comp. frequency (in events/s):	0	Component-6 Component-7	Host- Host-						Gree	edy Approxim	ation	
Maximum comp. frequency (in events/s):	10	Component-8							dice	заў Арргохії		
Minimum host reliability:	0	Component-9 Component-10								Clustered		
Maximum host reliability:	1	Component-11								Decentralize	_	
Minimum comp. event size (in KB):	0.01	Component-12								Decentralize	a	
Maximum comp. event size (in KB):	10	Component-13 Component-14										
Minimum host bandwidth (in KB/s):	30	Component-15					Numb	er of i	iterations:	1000		
Maximum host bandwidth (in KB/s):	1000	Component-16 Component-17										
Central host		Component-18			UseMapping		Bench	nmark	(how many t	imes): 1000		
Minimum bandwidth(in KB/s): 100		Component-19				-11	Ben	chmarl	kl			
Maximum bandwidth(in KB/s): 500		Component-20 Component-21										
Minimum reliability: .6		Component-22										
Maximum reliability: 1		Component-23										
		Component-24							Revert t	o previous de	eployment	
Generate		Component-25 Component-26										
		Component-27										
Availability:	0.81609	Component-28										
, tranazinty i		Component-29										
Tables of parameters				Results								
Hosts: reliability and memory Comps:	frequency and memory	Hosts: bandwidth		Component		Initial	I.d	E	Unbias	Biased	Greedy	Decent. 🔺
Hosts	n equerey and memory	, Lusses, sanamacu l	·	85		6			7	6	7	6
Hest/Hest 0 1	2 2	<u> </u>		86		Ē			4	7	7	7

Host/Host	0	1	2	3	4	5	
0	1.0	0.123	0.0	0.246	0.0	0.0	
1	0.123	1.0	0.0	0.0	0.0	0.0	
2 3	0.0	0.0	1.0	0.0	0.0	0.0	
3	0.246	0.0	0.0	1.0	0.684	0.0	
4	0.0	0.0	0.0	0.684	1.0	0.0	
5	0.0	0.0	0.0	0.0	0.0	1.0	
6	0.0	0.0	0.0	0.0	0.114	0.0	
7	0.672	0.883	0.627	0.966	0.630	0.781	
Mem	202.	235.	247.	232.	204.	247.	
•							

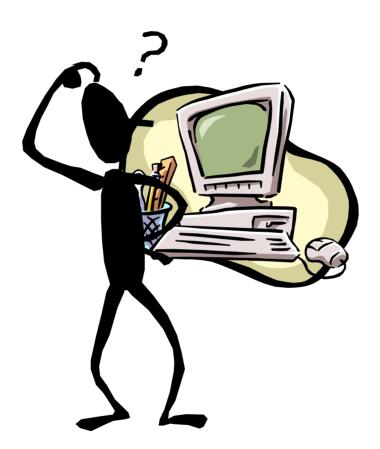
Component	Initial d	E	Unbias	Biased	Greedy	Decent.	٠
85	6		7	6	7	6	
86	6		4	7	7	7	
87	3		5	7	3	3	
88	1		0	7	1	6	
89	4		6	2	4	5	
90	7		6	7	7	7	
91	6		3	0	7	1	
92	6		1	7	6	1	
93	1		0	7	7	7	
94	1		6	6	6	3	
95	0		4	0	3	7	
96	0		3	4	4	0	
97	4		6	1	6	7	
98	0		5	6	3	5	
99	5		2	4	1	1	
Availability	0.3091		0.3937	0.4503	0.6334	0.6392.	
Running time (in ms)	0		90	601	7130	0	
Estimated redeployment time	N/A		20360	13149	16940	0.0	-



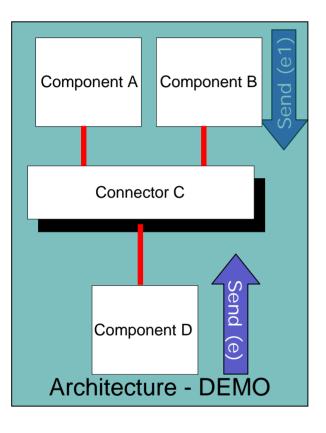




Questions?



Using Prism-MW



Component D sends an event

```
Event e = new Event ("Event_D");
e.addParameter("param_1", p1);
send (e);
```

Component B handles the event and sends a response

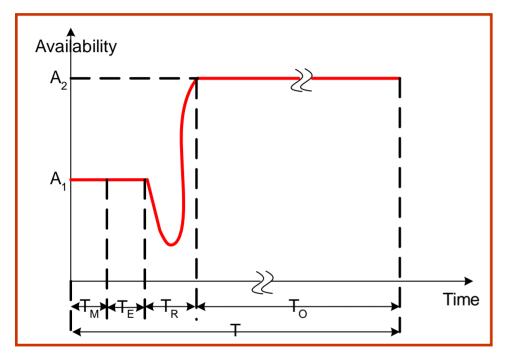
public void handle(Event e)

}

```
if (e.equals("Event_D")) {
```

```
Event e1= new Event("Response_to_D");
e1.addParameter("response", resp);
send(e1);
}...
```

Assumptions



 $T_M + T_F + T_R < < T_O$

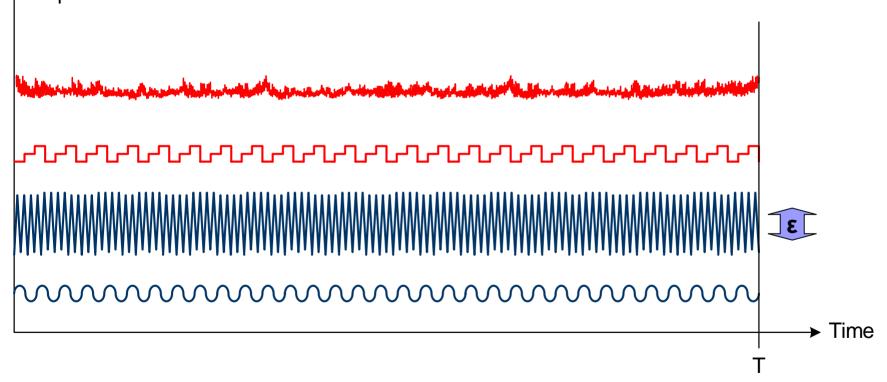
• *Possible in small amounts of time even for slow links*

For example,

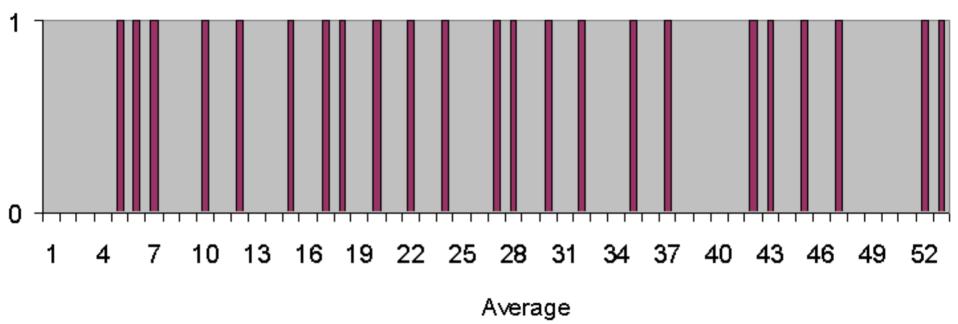
- dialup (56K), 50% reliability, 100K in 28 sec
- 1M, 50% reliability, 100K in less than 0.7 sec

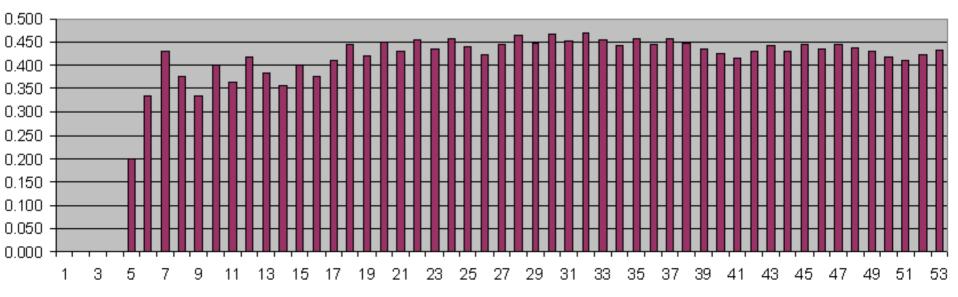
Assumptions

Average value of a parameter



Events





What Is Needed?

- Scalable and lightweight support for
 distributed architectures
 with arbitrary topologies
- Support for automatic monitoring
- Efficient and precise solutions for the exponentially complex redeployment problem
- Support for automatic redeployment

- 1) Lack of knowledge about runtime system parameters
 - unknown at the time of initial deployment
 - □ change at runtime
 - > architectural middleware with monitoring support

- 1) Lack of knowledge about runtime system parameters
 - unknown at the time of initial deployment
 - □ change at runtime
 - > architectural middleware with monitoring support

2) Exponentially complex problem

- \Box n components and k hosts = k^n possible deployments
- □ system constraints further complicate the problem
- > polynomial-time approximating algorithms

- 1) Lack of knowledge about runtime system parameters
 - unknown at the time of initial deployment
 - □ change at runtime
 - > architectural middleware with monitoring support

2) Exponentially complex problem

- \Box n components and k hosts = k^n possible deployments
- □ system constraints further complicate the problem
- > polynomial-time approximating algorithms

3) Assessing deployment architectures

- comparison of different algorithms
- performance vs. complexity, sensitivity analysis, etc.
- > analysis and simulation environment

- 1) Lack of knowledge about runtime system parameters
 - unknown at the time of initial deployment
 - □ change at runtime
 - > architectural middleware with monitoring support

2) Exponentially complex problem

- \Box n components and k hosts = k^n possible deployments
- □ system constraints further complicate the problem
- > polynomial-time approximating algorithms

3) Assessing deployment architectures

- comparison of different algorithms
- performance vs. complexity, sensitivity analysis, etc.
- > analysis and simulation environment

4) Effecting the selected solution

- redeploying components
- requires an automated solution
- > architectural middleware with deployment support

Why the Problem Isn't Solved

